

What is bandit problem

Mouse learning a T-maze:

The mice faced the dilemma of choosing to go left or right after starting in the bottom of a T-shaped maze, not knowing each time at which end they would find food. "Two-armed bandit" machine: Humans could choose to pull either the left or the right arm of the machine, each giving a random pay-off with the distribution of pay-offs for each arm unknown to the human player.('bandit' because they steal your money)

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Example: clinical trials

- You have multiple treatment options for cancer X.
- Cancer X patients arrive sequentially into your clinical trial.
- How do you treat them?
- Want to save as many lives as possible overall.

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Example: Feynman's Restaurant Problem



Example: A/B Testing

Whether the 'buy it now' button should be placed at the top of the product page or at the bottom.

- In traditional A/B testing, the objective of the statistician is to decide which website is better.
- When using a bandit algorithm, there is no need to end the trial. The algorithm automatically decides when one version of the site should be shown more often than another.

Key Elements of All of These Problems

We need to learn from past to predict the future and optimze. Core properties:

- · Sequentially taking actions of unknown quality
- The feedback provides information about quality of chosen action
- There is no state

It's a special tractable case of reinforcement learning.

Key Elements of All of These Problems

Key feature of the solution: balancing *exploitation* and *exploration*.

- Exploitation: use the current knowledge to focus on the action that seems to yield the highest rewards.
- Exploration: explore further the other actions to identify with better precision which action is actually the best.

Multi-armed bandit



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Multi-armed bandit









Mathematical Formulation

k arms, each with a reward distribution p_i . In each period t = 1, ..., T, we pull an arm $A_t \in \{1, ..., k\}$, and observe a reward $x_t \sim p_{A_t}$.

History: H_{t-1} = (A₁, x₁, A₂, x₂, ..., A_{t-1}, x_{t-1}).
Policy: π_t(H_{t-1}) = A_t

 $\min_{\pi} \mathbb{E}$

The objective is

or equvialently, to minimize the regret

 p_{A_t}

What is a good policy



Pure Greedy policy

Try each arm once. Then always pull the arm that appears the best (highest average reward to date). This can fail spectacularly:

• Consider a two-armed bandit: $x_1 \sim Ber(0.6)$, $x_2 \sim Ber(0.4)$.

- With probability 0.16, we start out by seeing 0 for arm 1 and 1 for arm 2.
- So we always pull arm 2 and accumulate 0.2 regret at each step.

• We have linear regret: Regret(T) $\geq cT$.

Explore then commit policy

Exploration phase: try each arm *m* times
Exploitation phase: for t = km + 1,..., n, pull the arm with the highest average reward.

1: Input m.

2: In round t choose action

$$A_t = egin{cases} (t mod k) + 1\,, & ext{if } t \leq mk\,, \ rgmax_i\,\hat{\mu}_i(mk)\,, & t > mk\,. \end{cases}$$

(ties in the argmax are broken arbitrarily)

where $\hat{\mu}_i(t) = \sum_{s=1}^t \mathbb{I}\{A_s = i\} x_s / \sum_{s=1}^t \mathbb{I}\{A_s = i\}$ denotes the empirical average reward of arm *i* up to time *t*.

Explore then commit policy



When ETC is interacting with any 1-subgaussian bandit and $1 \le m \le n/k$,

$$R_n \leq m \sum_{i=1}^k \Delta_i + (n-mk) \sum_{i=1}^k \Delta_i \exp\left(-\frac{m\Delta_i^2}{4}\right).$$

where $\Delta_i = \mu^* - \mu_i$ is the suboptimality gap between the mean of action *i* and the optimal action.

$$\sigma$$
-Subgaussian ($\mathbb{E}[X] = 0$): $E(exp(\lambda X)) \le \exp(-\sigma^2 \lambda^2/2), \forall \lambda$
Property: $\forall \epsilon > 0, \ P(X > \epsilon) \le \exp(-\epsilon^2/2\sigma^2).$

Concentration inequality

In probability theory, concentration inequalities provide bounds on how a random variable deviates from some value (typically, its expected value). Markov inequality (X nonnegative, a > 0):

 $\Pr(X \ge a) \le$

 $\mathsf{Pr}(|X - \mathsf{E}[X]| \ge a) \le \frac{\mathsf{Var}[X]}{\epsilon^2},$

 $\mathsf{P}(S_n - \operatorname{E}[S_n] \ge t) \le \exp\left(-\frac{2\tau}{\sum_{i=1}^n (b_i - a_i)^2}\right)$

• Chebyshev's inequality (E(X) and Var(X) finite):

• Hoeffding's inequality $(a_i \leq X_i \leq b_i)$:

Proof



Proof Assume without loss of generality that the first arm is optimal, which means that $\mu_1 = \mu^* = \max_i \mu_i$. By the decomposition given in Lemma 4.5, the regret can be written as

$$R_n = \sum_{i=1}^k \Delta_i \mathbb{E}\left[T_i(n)
ight] \,.$$

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(6.1)

In the first mk rounds, the policy is deterministic, choosing each action exactly m times. Subsequently it chooses a single action maximising the average reward during exploration. Thus,

$$\mathbb{E}[T_i(n)] = m + (n - mk)\mathbb{P}(A_{mk+1} = i)$$

$$\leq m + (n - mk)\mathbb{P}\left(\hat{\mu}_i(mk) \geq \max_{j \neq i} \hat{\mu}_j(mk)\right). \quad (6.2)$$

The probability on the right-hand side is bounded by

$$\begin{split} \mathbb{P}\left(\hat{\mu}_i(mk) \geq \max_{j \neq i} \hat{\mu}_j(mk)\right) &\leq \mathbb{P}\left(\hat{\mu}_i(mk) \geq \hat{\mu}_1(mk)\right) \\ &= \mathbb{P}\left(\hat{\mu}_i(mk) - \mu_i - (\hat{\mu}_1(mk) - \mu_1) \geq \Delta_i\right) \end{split}$$

The next step is to check that $\hat{\mu}_i(mk)-\mu_i-(\hat{\mu}_1(mk)-\mu_1)$ is $\sqrt{2/m}\text{-subgaussian},$

which by the properties of subgaussian random variables follows from the definitions of $(\hat{\mu}_j)_j$ and the algorithm. Hence by Corollary 5.5,

$$\mathbb{P}\left(\hat{\mu}_i(mk) - \mu_i - \hat{\mu}_1(mk) + \mu_1 \ge \Delta_i\right) \le \exp\left(-\frac{m\Delta_i^2}{4}\right). \tag{6.3}$$

Substituting Eq. (6.3) into Eq. (6.2) and the regret decomposition (Eq. (6.1)) gives the result. $\hfill \Box$



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MAB: Introduction

Explore then commit policy



 $R_n \leq m \sum_{i=1}^k \Delta_i + (n - mk) \sum_{i=1}^k \Delta_i \exp\left(-\frac{m\Delta_i^2}{4}\right).$ This illustrate exploration-exploitation tradeoff!

Explore too much (m large) then the first term is large.
Exploit too much (m small) then the second term is large.
If we choose m × n^{2/3}, then the regret R_n = O(Kn^{2/3}).
Note. the choice of m requires the knowledge of n.

ϵ -greedy policy

With probablity ϵ_t , explore an arm uniformly at random. With probablity $1 - \epsilon_t$, explore the arm with the highest reward so far. Again. ϵ_t trade-off between exploration and exploitation. If $\epsilon_t \simeq t^{-1/3} (K \log t)^{1/3}$, then the regret $R_t = O(t^{2/3}(K\log t)^{1/3})$ Similar regret rate to ETC But it does not require knowing n, the regret bound holds for any t (anytime algorithm) • the exploration is spread more evenly over time than ETC The regret of ETC and ϵ -Greedy aren't that great, because they do not adapt exploration to the observed data.

ϵ -greedy policy



- k=10
 A=1,2,3,4,5,6,7,8,9,10
- $\Pr{r|a} \sim N(q(a), 1)$

The optimal policy is to always choose a_3 , But we don't know that initially



ϵ -greedy policy





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MAB: Introduction

Upper Confidence Bound policy

Based on the idea of "optimism in the face of uncertainty." Algorithm:

- compute the empirical mean of each arm and a confidence interval;
- use the upper confidence bound as a proxy for goodness of arm.

Note: confidence interval chosen so that true mean is very unlikely to be outside of confidence interval

- 1: Input k and δ
- 2: for $t \in 1, \ldots, n$ do
- 3: Choose action $A_t = \operatorname{argmax}_i \operatorname{UCB}_i(t-1, \delta)$
- 4: Observe reward X_t and update upper confidence bounds

5: **end for**

UCB policy



UCB policy



Consider a UCB Algorithm on a stochastic k-armed 1-subgaussian bandit problem. For any horizon n, if $\delta = 1/n^2$, then

$$R_n \leq 3\sum_{i=1}^k \Delta_i + \sum_{i:\Delta_i>0} \frac{16\log n}{\Delta_i}.$$

Choosing $\delta \approx 1/n^2$, we have $R_n = O(\sqrt{Kn \log n})$. (Recall that ETC has $O(n^{2/3})$ regret.)

We still need the knowledge of n. In practice, we may use

$$UCB_i(t-1,\delta) = \hat{\mu}_i(t-1) + \alpha \sqrt{\frac{2\log t}{N_i(t-1)}}$$
 with a tuning parameter $\alpha > 0$.

UCB policy



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Bayesian bandit

- So far we have made no assumptions about the reward distribution R
- Bayesian bandits exploit prior knowledge of rewards p(R)
- They compute posterior distribution of rewards $p(R|H_t)$

 $BR_n =$

• Use posterior to guide exploration

 $(\mu_{A^*}-\mu_{A_t})$

Thompson Sampling

• Thompson Sampling is an alternative Bayesian method with similar regret guarantees.

- Before the game starts, the learner chooses a prior over a set of possible bandit environments. In each round, the learner samples an environment from the posterior and acts according to the optimal action in that environment.
- 1: Input Bayesian bandit environment $(\mathcal{E}, \mathfrak{B}(\mathcal{E}), Q, P)$

2: for
$$t = 1, 2, ..., n$$
 do

3: Sample
$$\nu_t \sim Q(\cdot | A_1, X_1, \dots, A_{t-1}, X_{t-1})$$

4: Choose $A_t = \operatorname{argmax}_{i \in [k]} \mu_i(\nu_t)$

5: **end for**

Note. the goal is to optimize Bayesian regret now.

Thompson Sampling



Let's consider a Beta distribution prior over the mean rewards of the Bernoulli bandits:

$$p(\theta_k) = \frac{\Gamma(\alpha_k + \beta_k)}{\Gamma(\alpha_k)\Gamma(\beta_k)} \theta_k^{\alpha_k - 1} (1 - \theta_k)^{\beta_k - 1}$$

$$\Gamma(n)=(n-1)!$$



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The posterior is also a Beta! Because beta is conjugate distribution for the Bernoulli distribution.

 $\frac{p(\theta \mid D)}{p(D)} = \frac{p(D \mid \theta) p(\theta)}{p(D)}$

A closed form solution for the bayesian update, possible only for conjugate distributions!

$$(\alpha_k, \beta_k) \leftarrow \begin{cases} (\alpha_k, \beta_k) & \text{if } x_t \neq k \\ (\alpha_k, \beta_k) + (r_t, 1 - r_t) & \text{if } x_t = k. \end{cases}$$

Thompson Sampling



Theorem 3.

Let $(\mathcal{E}, \mathcal{B}(\mathcal{E}), Q, P)$ be a k-armed Bayesian bandit environment such that for all $\nu \in \mathcal{E}$ and $i \in [k]$, the distribution $P_{\nu i}$ is 1-subgaussian (after centering) with mean in [0, 1]. Then the policy π of Thompson sampling satisfies

 $BR_n(\pi, Q) \leq C\sqrt{kn\log n}$



Common algorithms

ETC: $O(n^{2/3})$ regret	tanghy	tanghy
• ϵ -greedy: $O(t^{2/3})$ regr	et 🔊	(₂₀ ,
• UCB: $O(\sqrt{n \log n})$ regr	et at	ati
• Thompson sampling: C	$O(\sqrt{n \log n})$ regret	0
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Interactive bandits:		
https://pavlov.tech/203	19/03/02/animated-mult:	i-armed-bandit-policies/
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Other bandits

- Adversarial bandits: the learner must choose an action a_t and observe a reward r_t . The adversary chooses the reward distribution.
- **Contextual bandits**: the learner observes a context x_t at each time step. The learner must choose an action a_t and observe a reward r_t .

Non-stationary bandits: the reward distribution changes over time.

Reference

- Lattimore, T., & Szepesvári, C. (2020). Bandit Algorithms. *Cambridge University Press.* https://tor-lattimore.com/downloads/book/book.pdf
- Slivkins, A. (2019). Introduction to Multi-Armed Bandits. ArXiv. https://arxiv.org/pdf/1904.07272.pdf

MAB in dynamic pricing

Selling horizon: [0, T]Demand: Poisson process $\lambda_t = \lambda(p(t))$.

- Feasible prices: $[p, \bar{p}]$
- Inventory level: x

Algorithm:

• A "learning" phase (exploration) of length τ is used first, in which κ prices are tested

Then, a "pricing" phase (exploitation) fixes a "good" price based on demand observations in the first phase.

Omar Besbes, Assaf Zeevi, (2009) Dynamic Pricing Without Knowing the Demand Function: Risk Bounds and Near-Optimal Algorithms. Operations Research 57(6):1407-1420.

MAB in dynamic pricing

Algorithm 1. $\pi(\tau, \kappa)$ Step 1. Initialization: (a) Set the learning interval to be $[0, \tau]$, and the number of prices to experiment with to be κ . Put $\Delta = \tau/\kappa$. (b) Divide $[p, \overline{p}]$ into κ equally spaced intervals and let $\{p_i, i = 1, \dots, \kappa\}$ be the left endpoints of these intervals. Step 2. Learning/experimentation: (a) On the interval $[0, \tau]$ apply p_i from $t_{i-1} = (i-1)\Delta$ to $t_i = i\Delta$, $i = 1, 2, ..., \kappa$, as long as inventory is positive. If no more units are in stock, apply p_{∞} up until time T and STOP (b) Compute $\hat{d}(p_i) = \frac{\text{total demand over } [t_{i-1}, t_i)}{\Lambda}, \quad i = 1, \dots, \kappa.$ Step 3. Optimization: $\hat{p}^{u} = \arg\max\{p_{i}\hat{d}(p_{i})\},\$ Compute (9) $\hat{p}^c = \arg\min|\hat{d}(p_i) - x/T|,$ (10)and set $\hat{p} = \max\{\hat{p}^c, \hat{p}^u\}$. Step 4. Pricing:

Denote *n* as the market size: $x_n = nx$ and $\lambda_n(\cdot) = n\lambda(\cdot)$. Then, as $n \to \infty$, the regret is asymptotically optimal.

• Set $\tau \simeq n^{-1/4}$ and $\kappa \simeq n^{1/4}$ The regret bound locates in $[\frac{C'}{n^{1/2}}, \frac{C(\log n)^{1/2}}{n^{1/4}}]$

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Some of challenges



- Dynamic pricing
- Product assortment
- Advertising placement

Operational concerns:

- finite switches (limited number of prices to test))
- non-stationary reward (demand changes over time)
- switching cost (cost of switching prices)

People to follow

